

A Circle Method Approach to K-Multimagic Squares

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What is a K -multimagic square

Definition

Given $K \geq 2$ we say a matrix $\mathbf{Z} \in \mathbb{Z}^{N \times N}$ is a K -multimagic square of order N or **MMS**(K, N) for short if the matrices

$$\mathbf{Z}^{\circ k} := (z_{i,j}^k)_{1 \leq i,j \leq N},$$

remain magic squares for $1 \leq k \leq K$.

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Definition

If a **MMS**(K, N) contains every integer from 1 to N^2 then it is called a *normal* **MMS**(K, N).

Trivial Examples of $\mathbf{MMS}(K, N)$

1	1	1
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This is a $\mathbf{MMS}(K, 3)$ for all
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a	b	c	d
d	c	b	a
b	a	d	c
c	d	a	b

For any $a, b, c, d \in \mathbb{Z}$ this is a $\mathbf{MMS}(K, 4)$ for all $K \geq 2$.

Definition of trivial $\mathbf{MMS}(K, N)$

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Question: Given $K \geq 2$ and $N \geq 4$ does there exist nontrivial $\mathbf{MMS}(K, N)$?

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17	36	55	124	62	114
58	40	129	50	111	20
108	135	34	44	38	49
87	98	92	102	1	28
116	25	86	7	96	78
22	74	12	81	100	119

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Open Problem

Does there exist a nontrivial $\mathbf{MMS}(2, 5)$?

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1	22	33	41	62	66	79	83	104	112	123	144
9	119	45	115	107	93	52	38	30	100	26	136
75	141	35	48	57	14	131	88	97	110	4	70
74	8	106	49	12	43	102	133	96	39	137	71
140	101	124	42	60	37	108	85	103	21	44	5
122	76	142	86	67	126	19	78	59	3	69	23
55	27	95	135	130	89	56	15	10	50	118	90
132	117	68	91	11	99	46	134	54	77	28	13
73	64	2	121	109	32	113	36	24	143	81	72
58	98	84	116	138	16	129	7	29	61	47	87
80	34	105	6	92	127	18	53	139	40	111	65
51	63	31	20	25	128	17	120	125	114	82	94

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74	8	106	49	12	43	102	133	96	39	137	71
140	101	124	42	60	37	108	85	103	21	44	5
122	76	142	86	67	126	19	78	59	3	69	23
55	27	95	135	130	89	56	15	10	50	118	90
132	117	68	91	11	99	46	134	54	77	28	13
73	64	2	121	109	32	113	36	24	143	81	72
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Open Problem

Does there exist a nontrivial $\text{MMS}(3, 11)$?

Formulating our problem

Given $K \geq 2$ let $N(K)$ denote the smallest natural number for which there exists non-trivial $\mathbf{MMS}(K, N(K))$.

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K	Upper bound on $N(K)$	Attributed to
2	6	J. Wroblewski
3	12	W. Trump
4	243	P. Fengchu
5	729	L. Wen
6	4096	P. Fengchu
$K \geq 2$	$(4K - 2)^K$	Zhang, Chen, and Lei

Our results

Via the Hardy-Littlewood circle method we establish

Theorem (F. 2024+)

$N(K) \leq 2K(K + 1) + 1$ for $K \geq 2$.

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Theorem (F. 2024+)

$$N(K) \leq 2K(K + 1) + 1 \text{ for } K \geq 2.$$

- ▶ This beats previously known results as soon as $K \geq 4$ and shows that $N(K)$ grows at most quadratically in K rather than potentially exponential in K .

Our results

- ▶ One may prove an analogous statement for prime valued $\mathbf{MMS}(K, N)$ by reapplying the entirety of the circle method where we detect prime solutions instead of integer solutions.

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- ▶ This, however, is not necessary as via an argument due to Granville in his expository article *Prime Number Patterns* one may apply the Green-Tao theorem and deduce the following.

Corollary (F. 2024+)

Given $K \geq 2$ there exists infinitely many nontrivial prime valued $\mathbf{MMS}(K, N)$ for every $N > 2K(K + 1)$.

Brief overview of our methods

Let $C = (c_{i,j})_{\substack{1 \leq i \leq r \\ 1 \leq j \leq s}} \in \mathbb{Z}^{r \times s}$ be given, and consider the diagonal system

$$\sum_{1 \leq j \leq s} c_{i,j} x_j^k = 0 \quad (1 \leq i \leq r, \quad 1 \leq k \leq K). \quad (1.1)$$

We define $R_K(P; C)$ to be the number of solutions $\mathbf{x} \in \mathbb{Z}^s$ to (1.1) where $\max_j |x_j| \leq P$.

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There exists of a matrix $C_N^{\text{magic}} \in \{-1, 0, 1\}^{2N \times N^2}$ for which $R_K(P; C_N^{\text{magic}})$ counts the number of **MMS**(K, N) with entries satisfying

$$\max_{1 \leq i, j \leq N} |z_{i,j}| \leq P.$$

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The number of trivial **MMS**(K, N) counted by $R_K(P; C_N^{\text{magic}})$ is at most $O(P^N)$, thus if one wishes to establish the existence of non-trivial **MMS**(K, N) it is enough to show that

$$\frac{R_K(P; C_N^{\text{magic}})}{P^N} \rightarrow \infty \quad \text{as} \quad N \rightarrow \infty. \quad (1.2)$$

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Theorem (F. 2024+)

For $K \geq 2$ and $N > 2K(K+1)$ there exists a constant $c > 0$ for which one has the asymptotic formula

$$R_K(P; C_N^{\text{magic}}) \sim cP^{N(N-K(K+1))}.$$

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However, in their application of the circle method they require the $r \times s$ matrix of coefficients C to be *highly non-singular*, i.e., for all $J \subset \{1, \dots, s\}$ with $|J| = r$ one should have

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Our matrix of interest C_N^{magic} does not satisfy this property. Thus we developed a version of the circle method for this problem that works if C satisfies some weaker property.

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For a given $r \times s$ matrix $C = [\mathbf{c}_1, \dots, \mathbf{c}_s]$ and any set $J \subset \{1, \dots, s\}$, we denote by C_J the submatrix of C consisting of the columns \mathbf{c}_j where $j \in J$.

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For any $a \in \mathbb{Z}$ and $b \in \mathbb{N}$ we denote by $\text{rem}(a, b)$ the remainder of a modulo b considered as an integer between 0 and $b - 1$.

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Definition

We say that a matrix C dominates a function $f : \mathbb{N} \rightarrow \mathbb{R}^+$ whenever the inequality

$$\text{rank}(C_J) \geq \min \{f(|J|), r\},$$

holds for all $J \subset \{1, \dots, s\}$.

Brief overview of our methods

Quantitative Hasse Principle (F. 2024+)

Let $K \geq 2$ and suppose that $C \in \mathbb{Z}^{r \times s}$ satisfies $s > rK(K+1)$.
Then, if C dominates the function

$$F(x) = \max \left\{ \frac{x - \text{rem}(s, r)}{\lfloor \frac{s}{r} \rfloor}, \frac{x - \text{rem}(s-1, r)}{\lfloor \frac{s-1}{r} \rfloor} \right\},$$

one has that

$$R_K(P; C) = P^{s - \frac{rK(K+1)}{2}} (\sigma_K(C) + o(1)),$$

where $\sigma_K(C) \geq 0$ is a real number depending only on K and C .
Additionally $\sigma_K(C) > 0$ if there exists non-singular real and p -adic solutions to the system (1.1).

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- ▶ The matrix C_N^{magic} dominates the the function $F(x)$ defined previously.
- ▶ This is done via a combinatorial argument and understanding the underlying linear system associated to the matrix C_N^{magic} .
- ▶ Establish the existence of nonsingular real and p -adic **MMS**(K, N).
- ▶ This is done by looking at a particular integer valued **MMS**(K, N) and showing that the Jacobian of the associated linear system at that that point is full rank.

Breaking news!

[Submitted on 13 Jun 2024]

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Theorem 1.3. *Let $d \geq 3$. There is a positive integer $n_0(d)$ such that for every integer $n \geq n_0(d)$, there exists an $n \times n$ magic square of d^{th} powers.*

$$n_0(d) = \begin{cases} 14 \min\{2^d, d(d+1)\} + 79 & \text{if } 3 \leq d \leq 19, \\ 14 \lceil d(\log d + 4.20032) \rceil + 79 & \text{if } d \geq 20. \end{cases}$$

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What would our methods say if applied to this single degree case?

Single degree case

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Confident can show

For $d \geq 3$ there exists a nontrivial $N \times N$ square of d th powers for all $N \geq d(\log(d) + 4.20032) + 1$.

Thank you for listening!